

## Lecture 22-1348

### Note on Review Session

Options	Wednesday, April 16 <sup>th</sup>	1 PM	*
	Thursday April 17 <sup>th</sup>	1 PM	
	Friday April 18 <sup>th</sup>	1 PM	

Last Time 9.3 Graph Isomorphism Note It will be outlined.

This time 9.4. let  $G$  be a graph. So  $G = (V, E)$

A path in  $G$  is a sequence of vertices

$$v_0, v_1, \dots, v_n$$

s.t. each  $v_i$  is adjacent to  $v_{i+1}$

We allow paths to have repeated vertices, i.e.



etc.

A path is a cycle if it begins & ends at same vertex



~~Definition~~ A path is simple if it has no repeated vertices.

A circuit is simple if it has no repeated vertices except starting/ending point.

Thm: In a graph  $G$ , say that two vertices  $a$  and  $b$  are related,  $a \sim b$ , if there is a path from  $a$  to  $b$ .

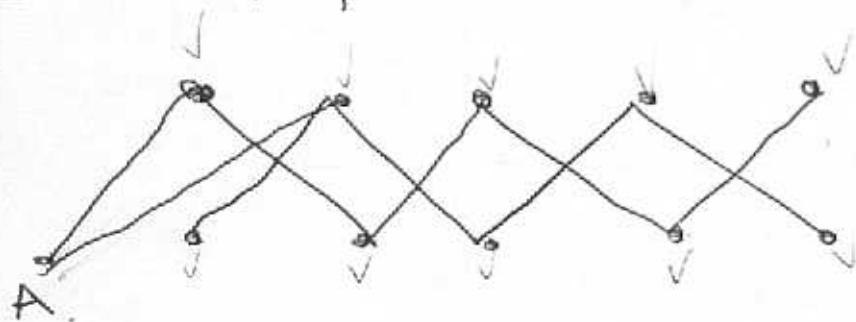
This is an equivalence relation.

Proof last time. Know it.

Defn: A graph is connected if between any 2 vertices there is a path.

Ex: #4, 5, p 630

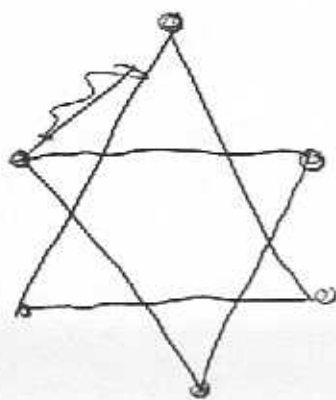
#4



A.

Note from point A, I can get to every point. Therefore there is a path between any 2 points. Why?

#5



Not connected

Thm: A graph is connected if and only if  
for all vertices  $a \in V$

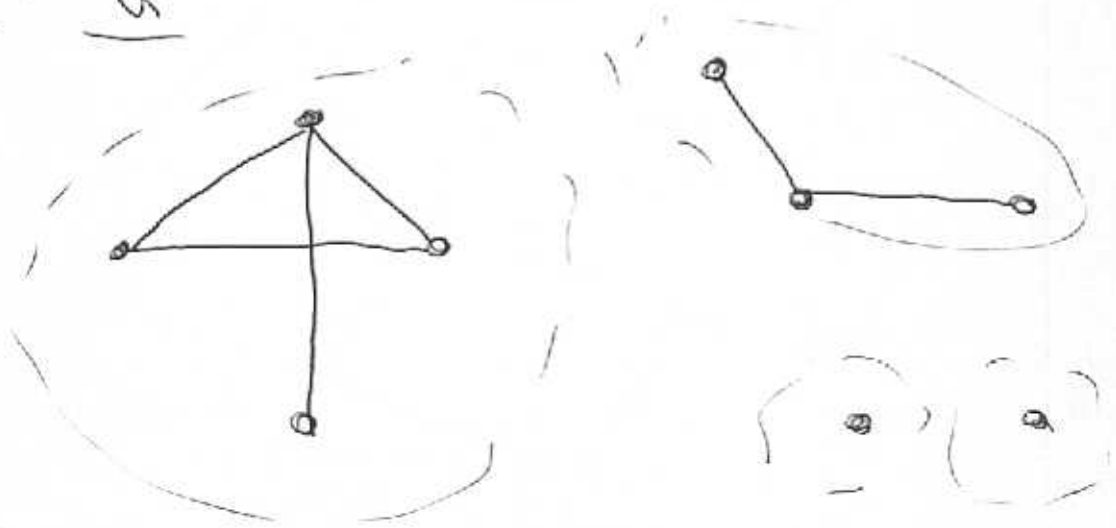
$$[a] = \text{equivalence class of } a = V$$

Note Suppose the graph  $G$  is not connected.

Then if  $a \in V$ ,  $[a] \subseteq V$ , but is not all of  $V$

~~All~~ All those vertices which are in  $[a]$   
form the connected component of  $a$ .

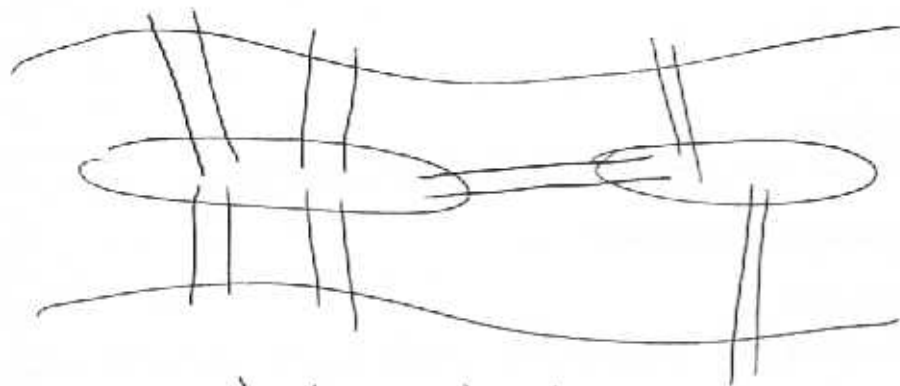
Ex:  $G$



4 connected components.

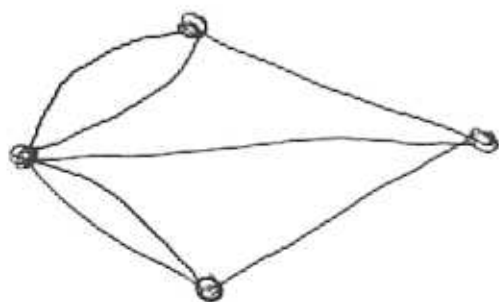
~~Section~~ Section 9.5, last section

Königsberg bridge system



Is there a path through Königsberg which crosses each bridge exactly once and returns to starting point? You can start wherever you want.

This can be drawn more abstractly as



Note This is a multigraph, not a graph.

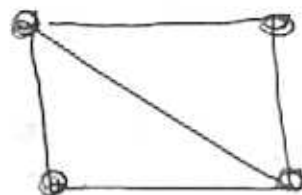
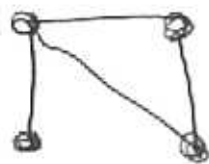
~~We can~~ Def'n p 633

Let  $G$  be a connected graph. An Euler circuit is a simple circuit (remember what this means) which contains every ~~vertex~~ edge exactly once.

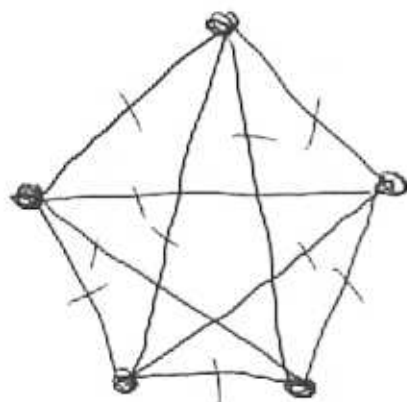
Note, in particular, an EC starts and ends at same point, since it is a circuit.

Q: Must a connected graph have a circuit?

No Ex



Both fail. Here's one that works



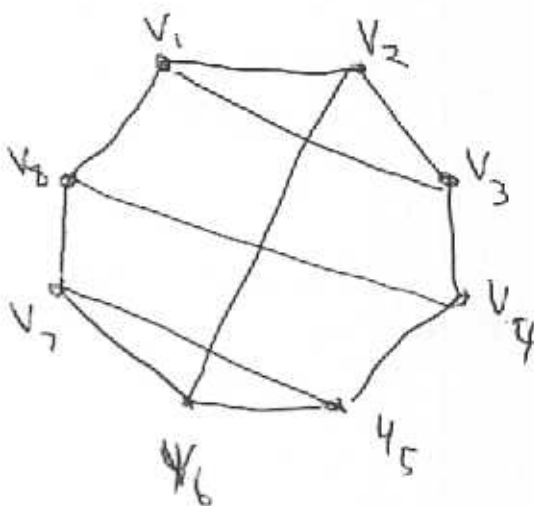
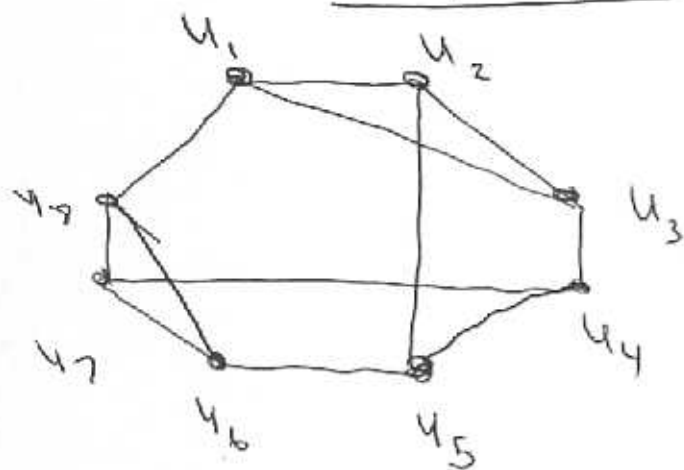
This is  $K_5$

Number edges.

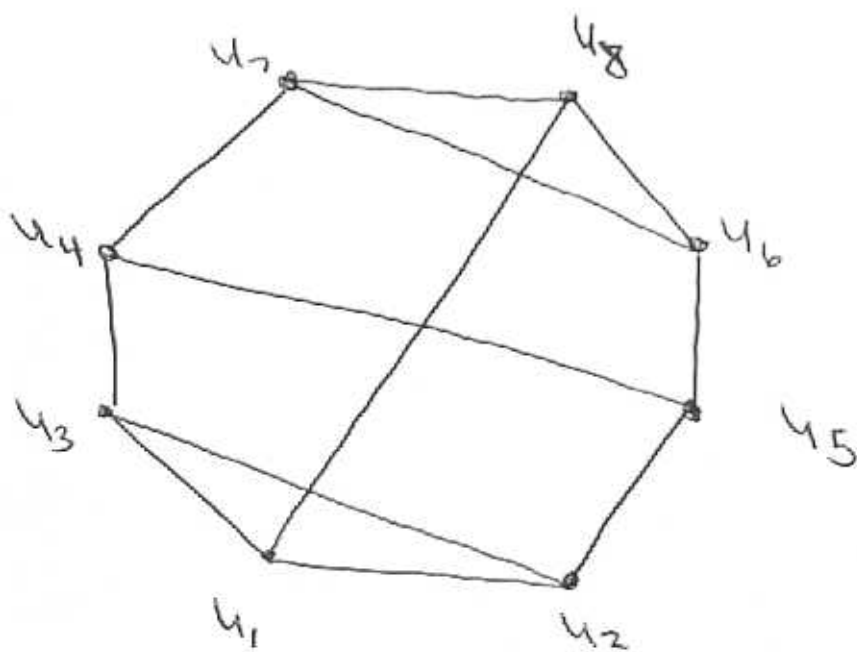
What is the theorem?

# Graph Iso Example

#20, p631



Claim - They are iso.



Add names in this order  $u_8, u_1$  first  
 Then  $v_7, u_6$  (can go on either side)  
 Then  $u_3, u_2$  (this is forced)  
 $u_4$  is forced  
 $u_5$  is forced.